

# Data structures in Information Retrieval

**Basic structures and environments**

# Agenda

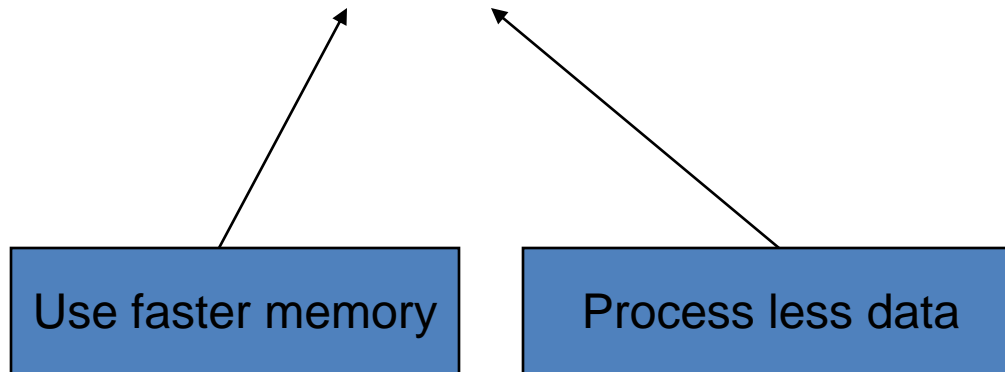
- How to create an effective structure
- Modern hardware
- Compression
- Search structures
- Parallelization

# General approach

1. Think about data usage
2. Select algorithm (s)
3. Select structure (s)
4. Select right libraries

# Effective data representation

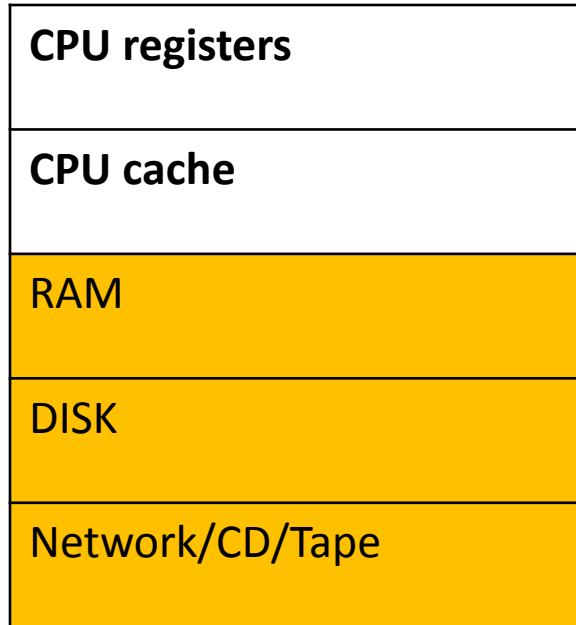
1. Less memory resources
2. Faster execution



# Modern Hardware

- Moore law
- Multilayer memory
- High parallelism

# Memory hierarchy



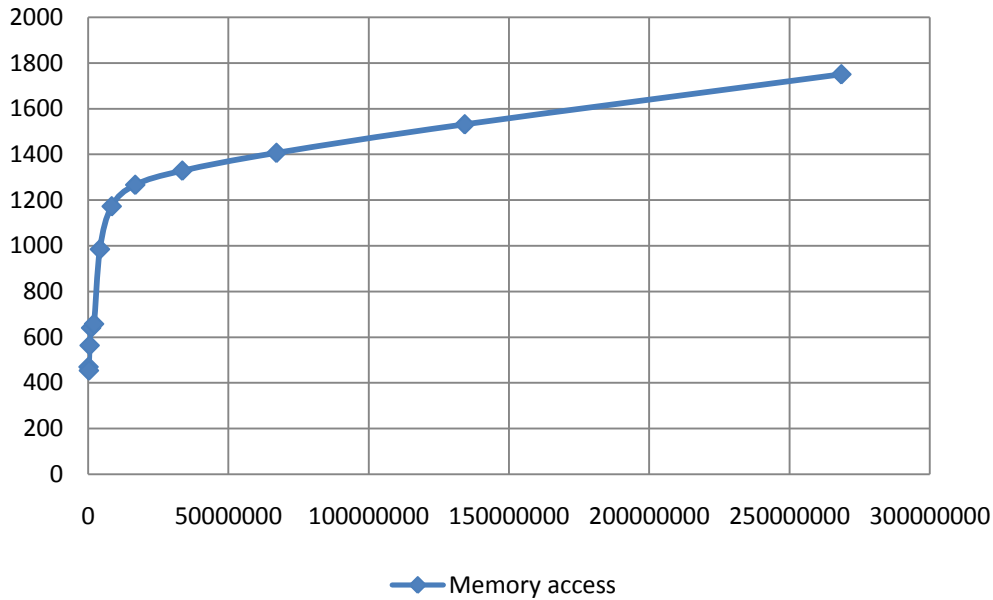
NUMA



# RAM access

```
unsigned val = 0;
for(unsigned k = 0;k<1024*1024*8;++i){
    pos = rand()%memSize;
    unsigned t = arr[pos];
    arr[pos] = val;
    val = t;
}
```

## Memory access

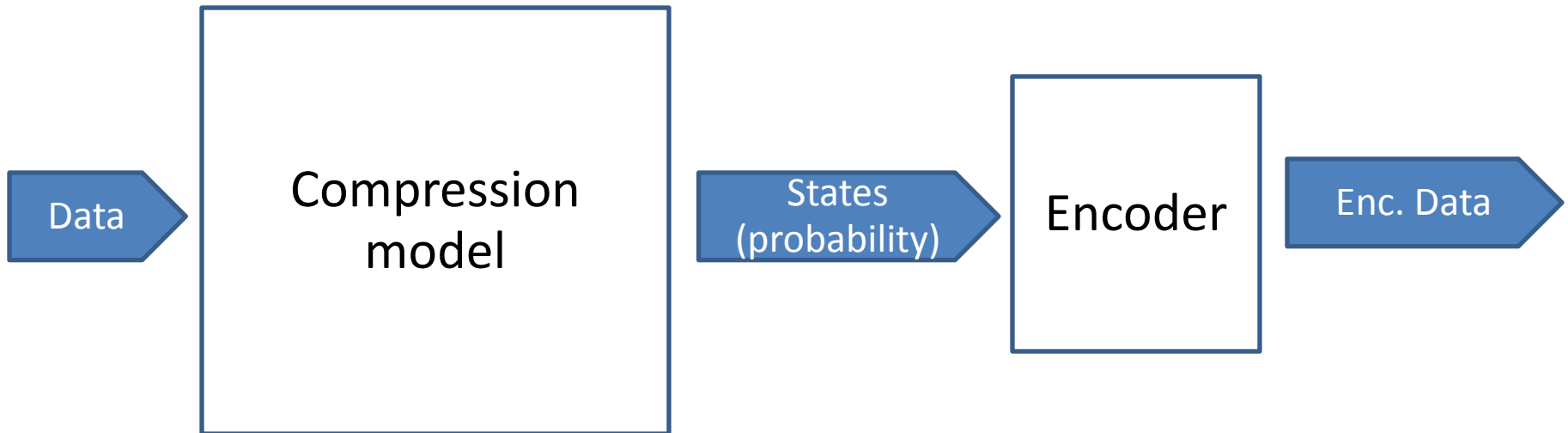


# Storing data

- Data (documents): text + markup
- Index structures: words, positions, flags



# Data encoding and compression



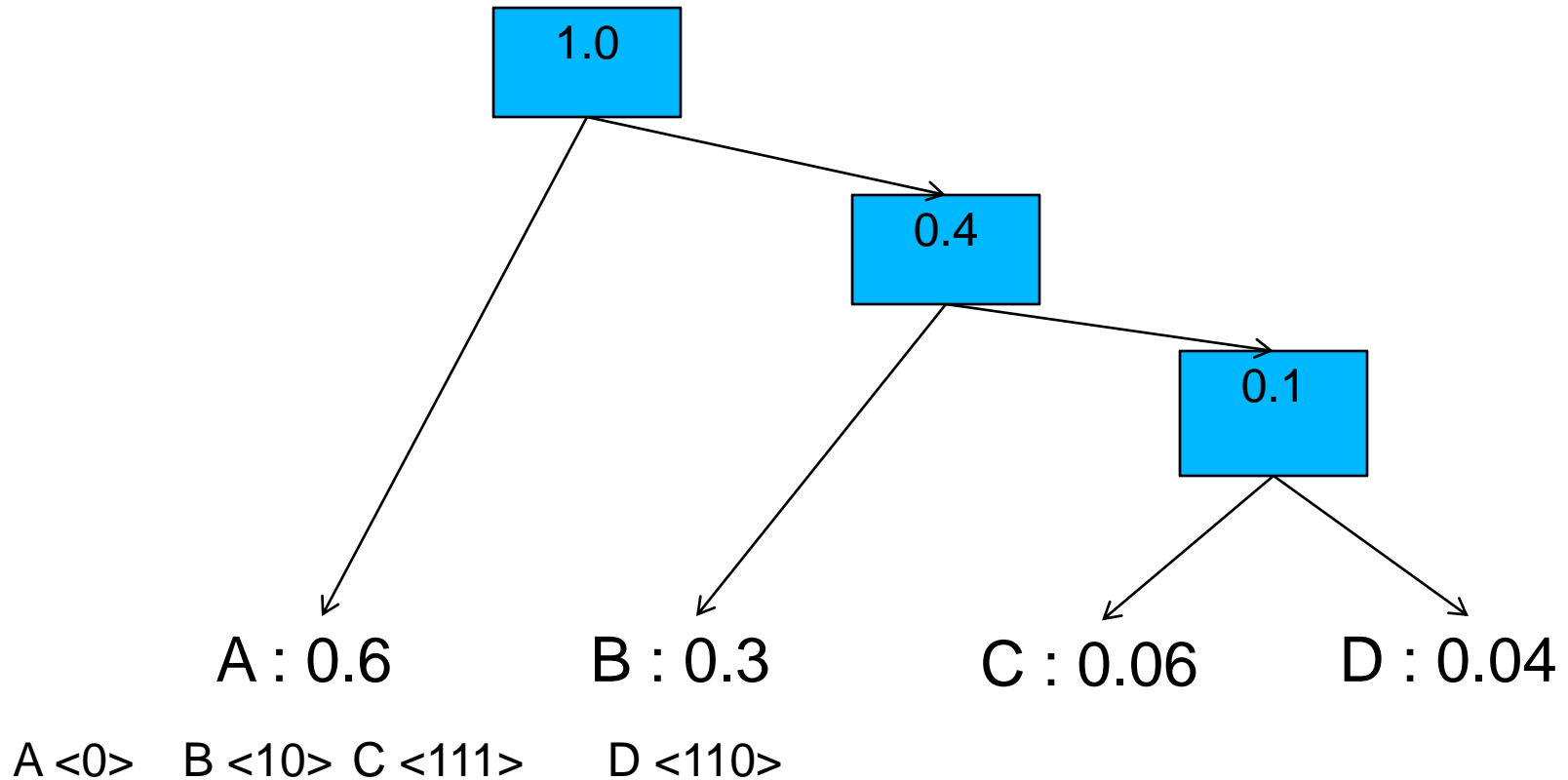
# Entropy coders

Input: Codes from alphabet  $A$

probabilities distribution  $a \rightarrow P_a$

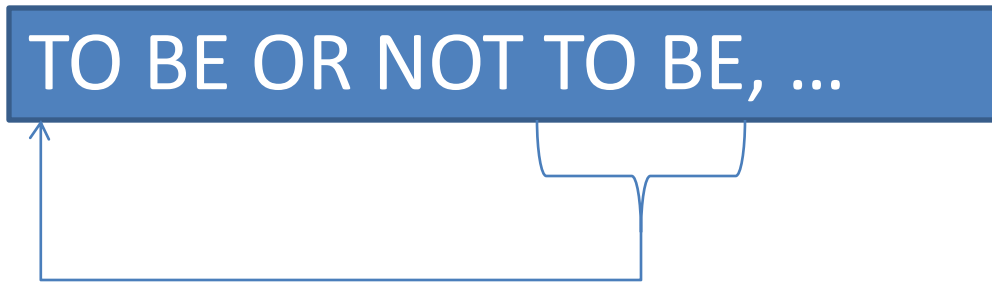
Output: minimal length encoding (minimal number of bits)

# Huffman coding



# LZ encoding

Add special code “back reference”



'T','O',' ','B','E',' ','O','R',' ','N','O','T',' ','ESC,8,4

# Storing text for “toy” search Engine

- Use compression (read + LZW decompression)
- Compressed index+data < data

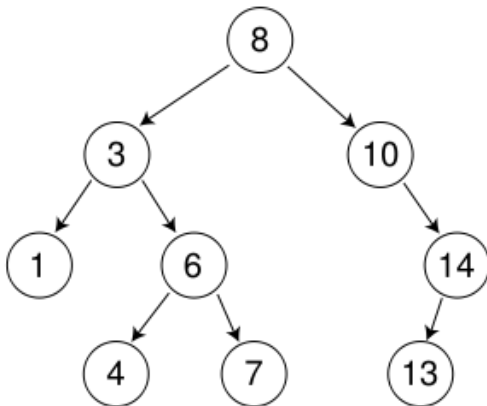
# Storing other data

Flags, field codes, etc. – entropy coder:

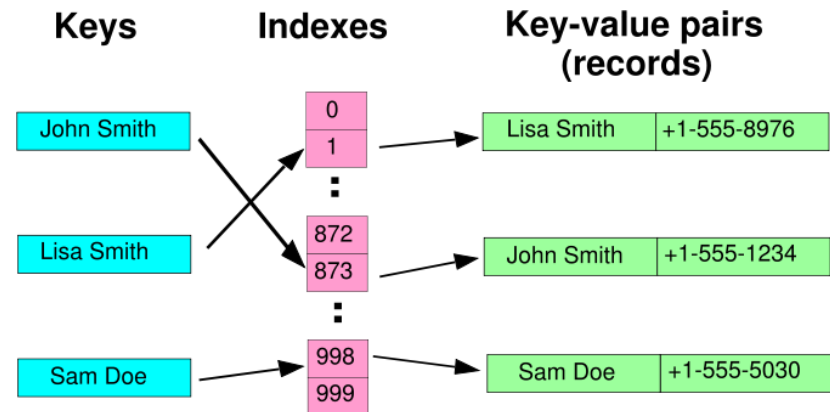
1. Collect statistics of this values
2. Select encoding
3. Store encoded in data and index

# Search structures

Branch & Bounds



Mappings



# Sorted arrays

A
...
BALL
BAT
...
MAN
MEN
...
ZAP
ZOO
ZZZ

Search:  $O(\log(N))$

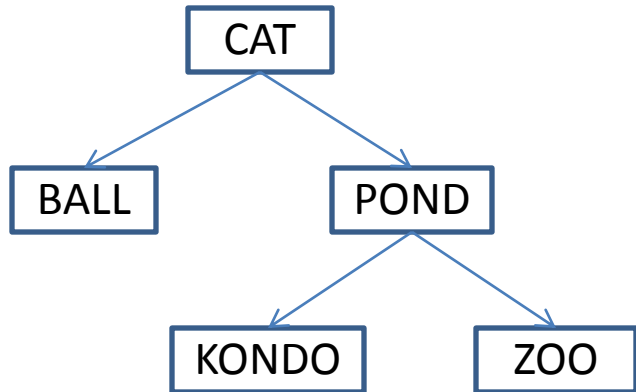
Insert:  $O(N)$

Advantage: most compact  
(no pointers)

Perfect for static dictionaries



# Binary tree

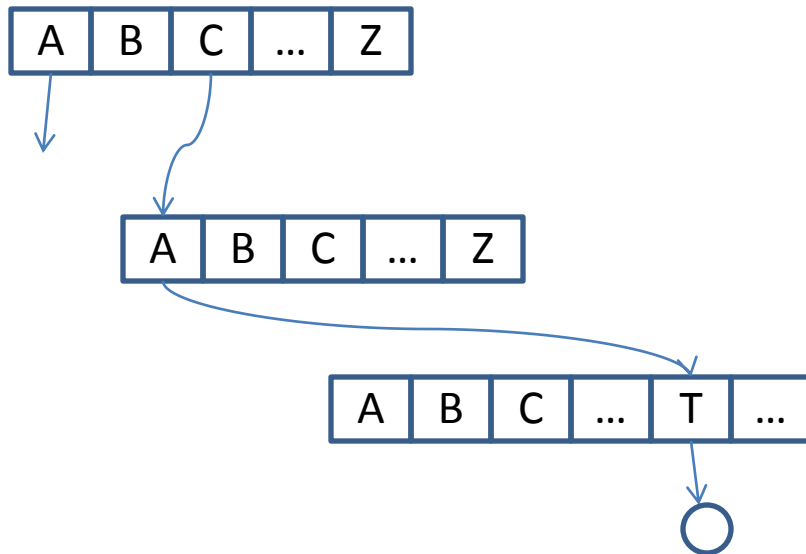


Search:  $O(\log(N))$

Insert:  $O(\log(N))$

Problems: pointers, memory allocations

# Trie

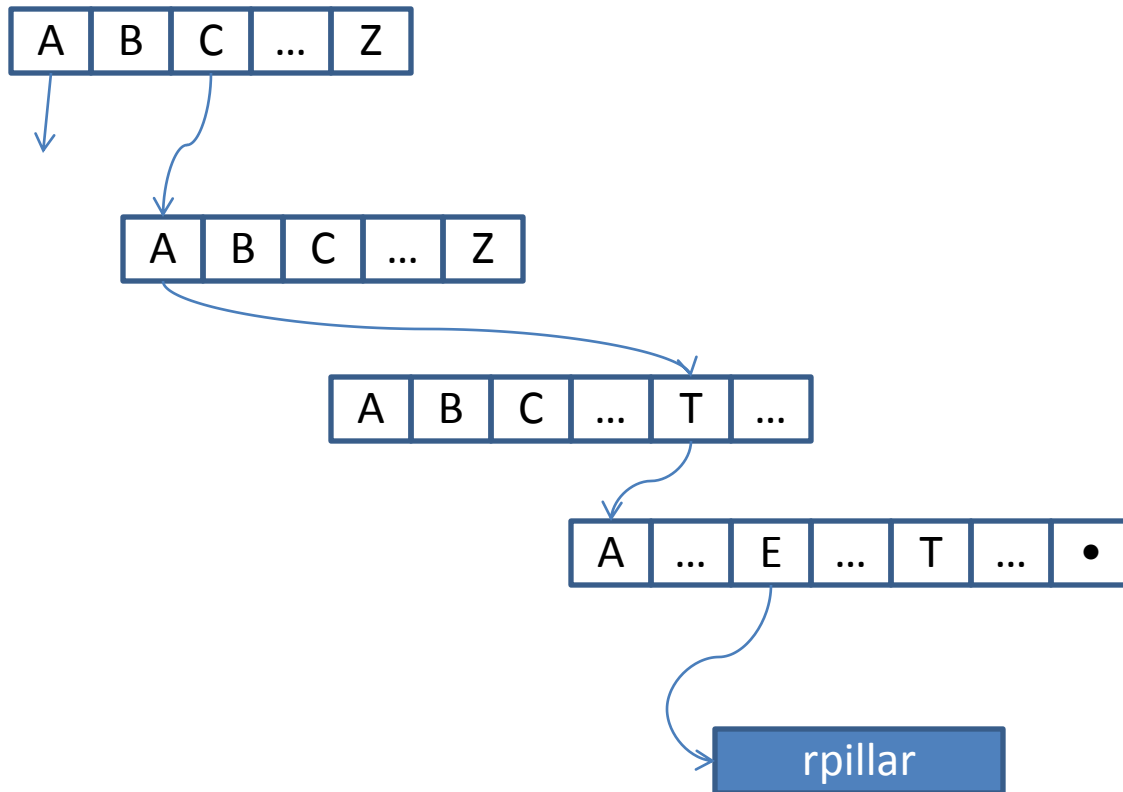


Search:  $O(Q)$

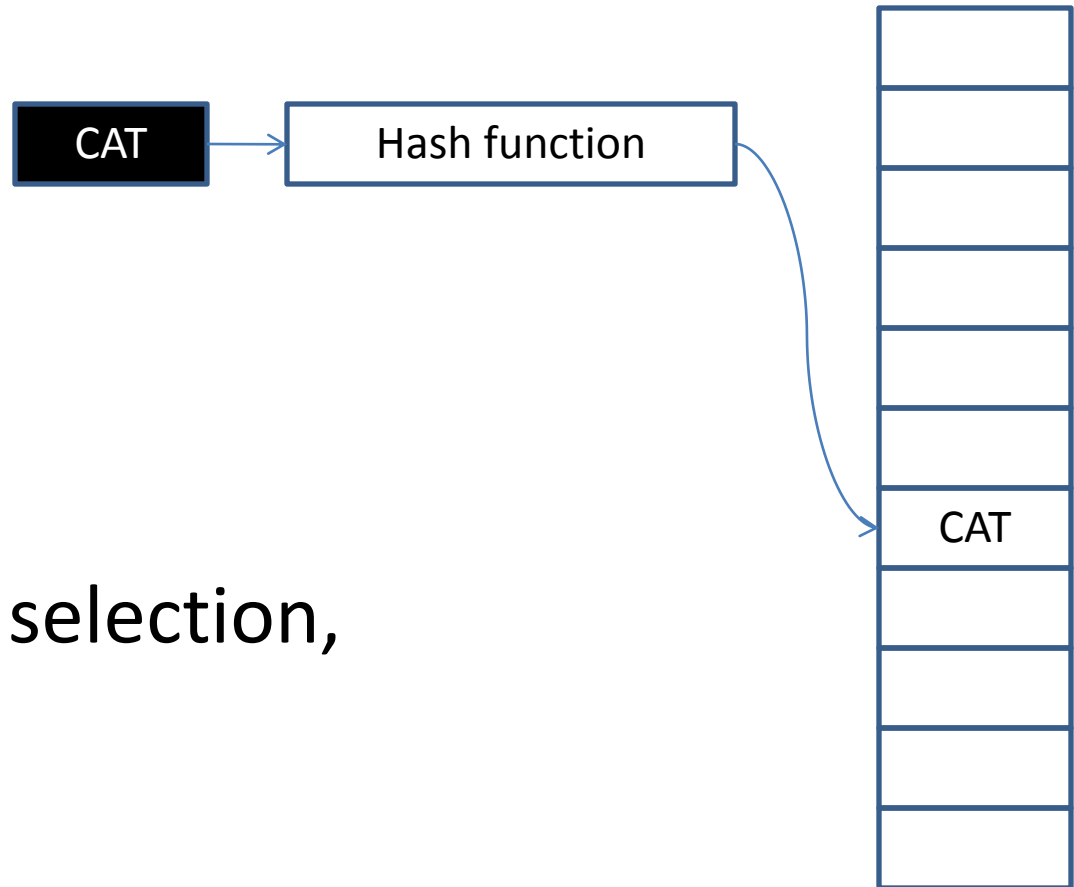
Insert:  $O(Q)$

Problems: memory

# Compressed trie



# Hash

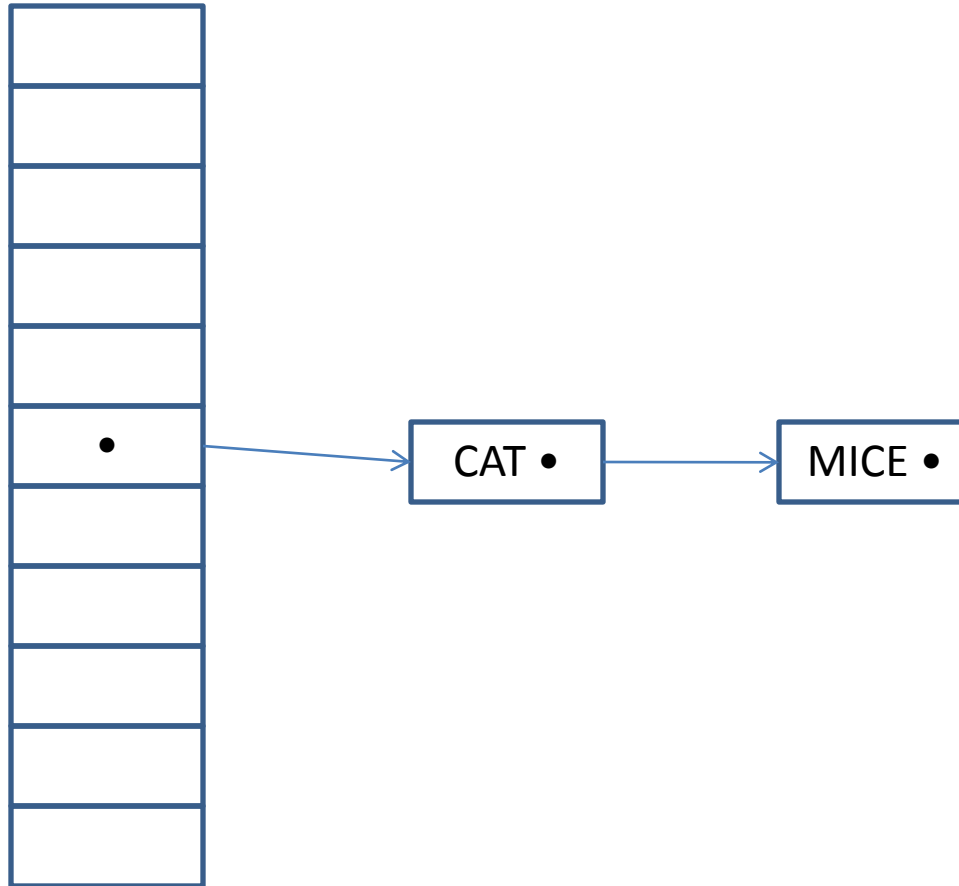


Search:  $O(1)$

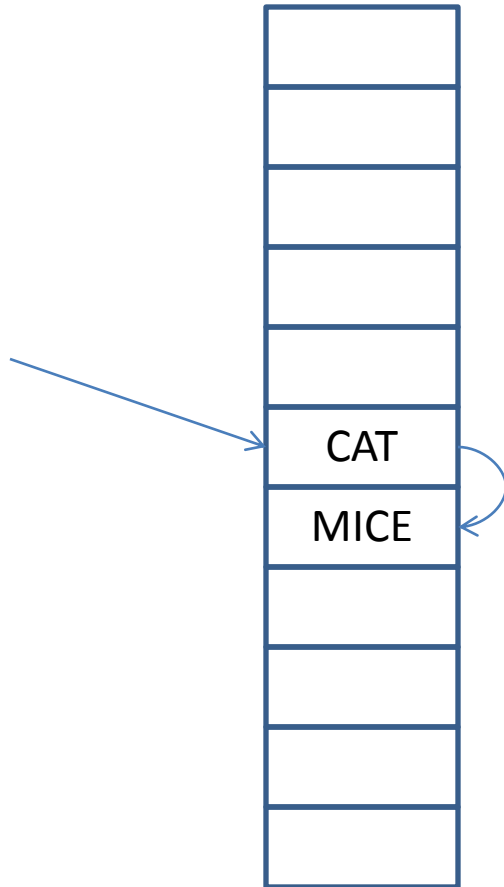
Insert:  $O(1)$

Problems: function selection,  
size of array

# Hash overflow [lists]



# Hash overflow [rehashing]



# Bloom filter

Search-acceleration structure:

0 – if key doesn't exist

1 – there is probability  $P$  that key exists

# Example: speller

## Problem:

- User input a word
- Suggest possible spell corrections

## Data Analysis:

Generate correction modification rules

$a \rightarrow aa(0.000005); l \rightarrow k(0.0006); ea \rightarrow ee(0.0004)$

## Algorithm:

Apply up to 2 rules to a word and return variants that have higher probability



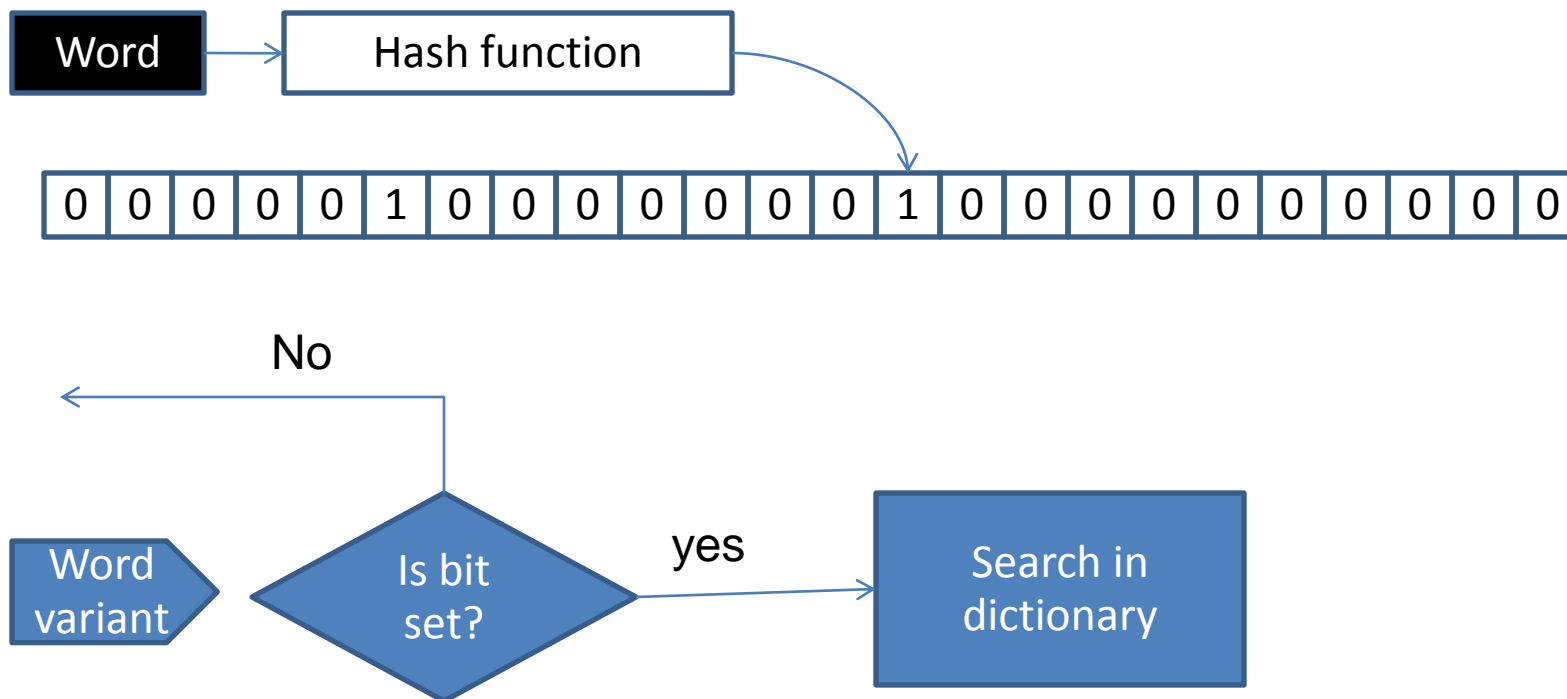
# Speller problem

For every word we can generate mlns variants,  
most of them – **not-valid words**

Search in big dictionary – relatively slow.

Idea – built a bloom filter that accelerates  
search

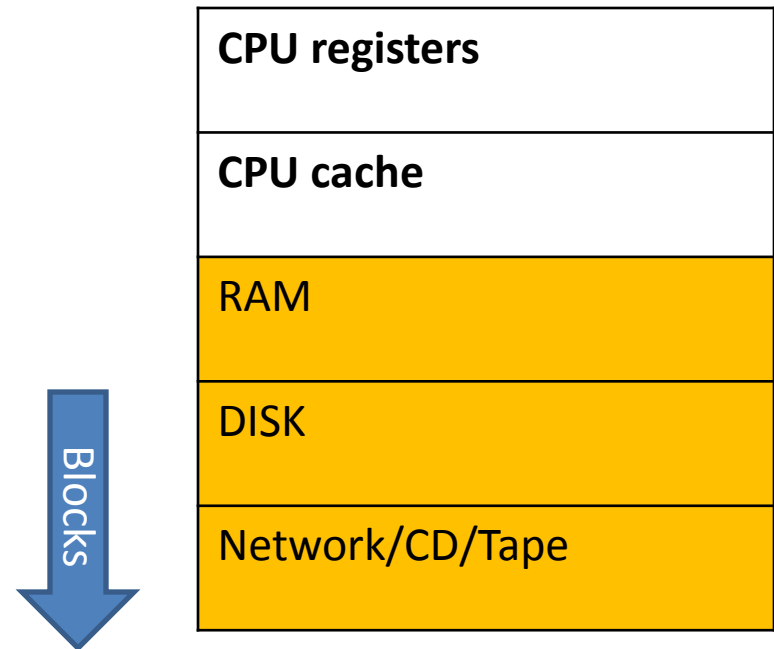
# Bloom filter for speller



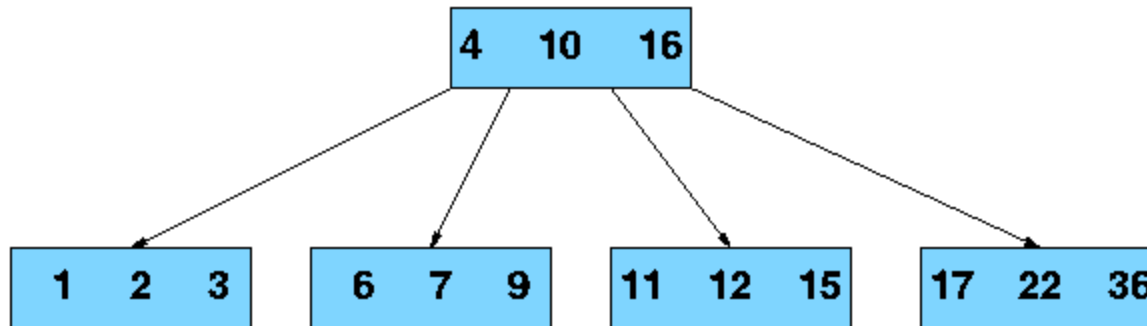
Increase speed of speller analysis 100 times

# Secondary-storage structures

- Block access
- Persistence
- Slow random access

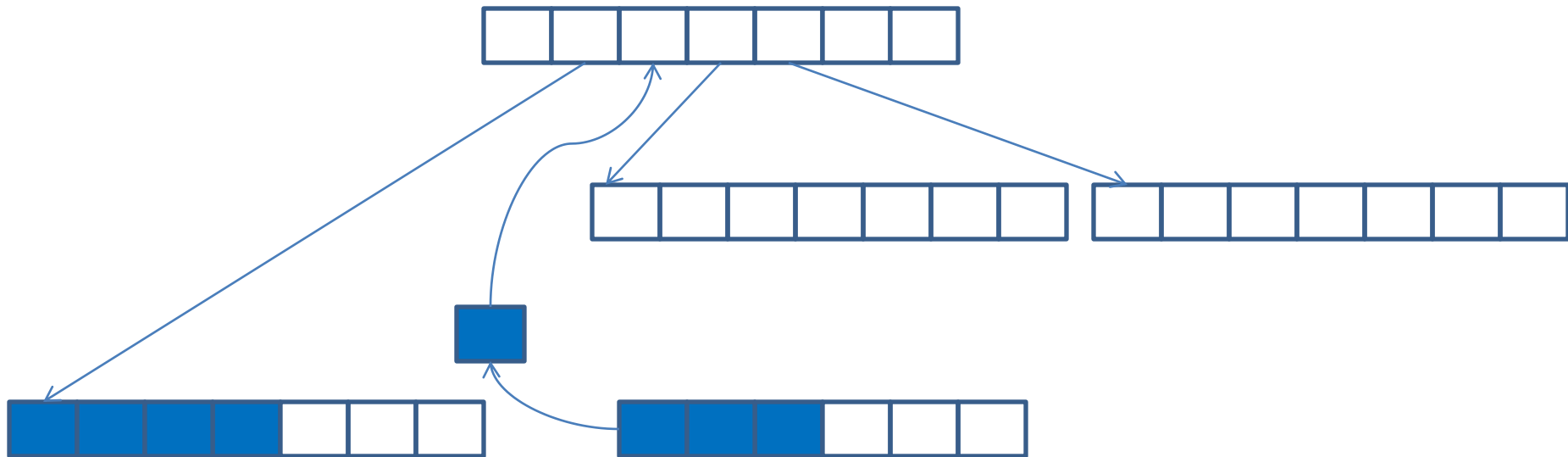


# B-Tree



- less pointers
- prefix-compressed nodes

# Insert in B-Tree



# Scaling your application

- **Vertical scaling** – moving to larger computer
- **Horizontal scaling** – dividing application into parts:

Two dimensions:

- **Grouping data by functions**
- **Grouping functions by data**

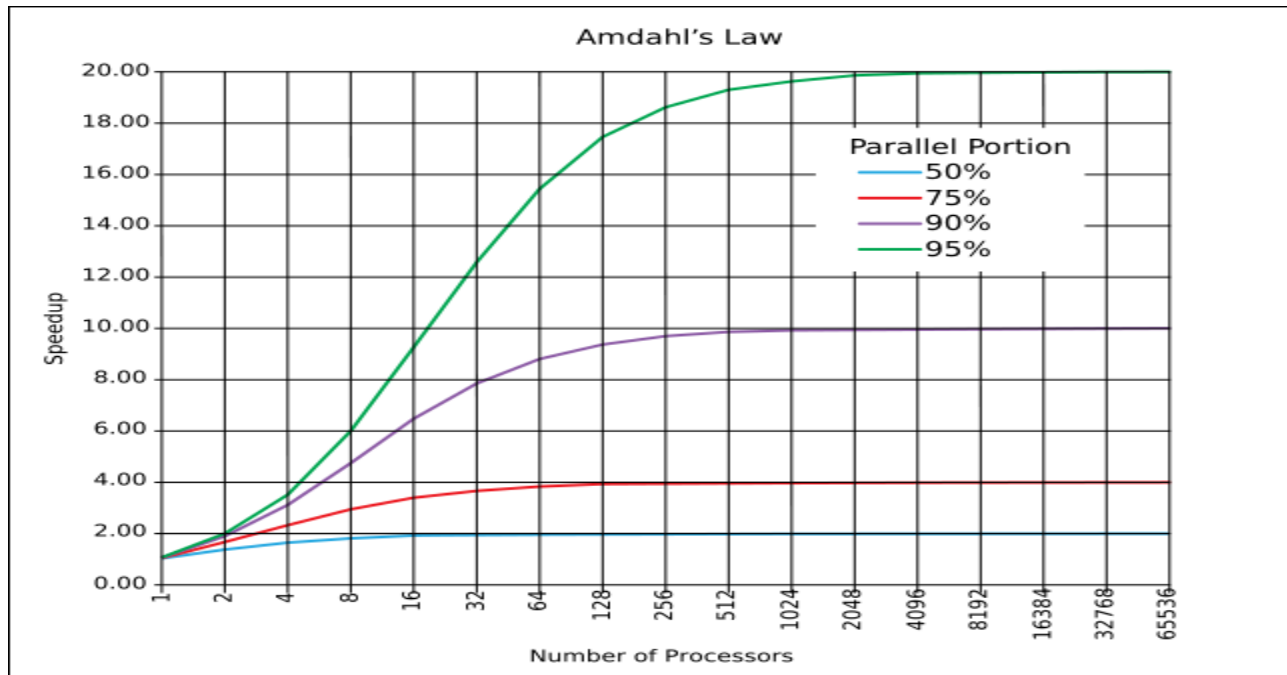
# Why scalability is not linear?

- Communication/Synchronization/Bookkeeping overhead
- Not all resources can grow linear

# Simple scalability model: Amdahl's law

$$\frac{1}{(1 - P) + \frac{P}{N}}$$

P - % of parallelized part, N – number of unit



RuSSIR

Russian Summer School  
in Information Retrieval

2008



# Parallel Execution

Serial Computing Is Dead; the Future Is  
Parallelism

**SearchDataCenter.com (06/30/08)**

“Sequential programming is really hard, and  
parallel programming is a step beyond that.”

**Andrew S.**

# Parallel architectures

- Multi CPU (Shared memory)
- Cluster of computers

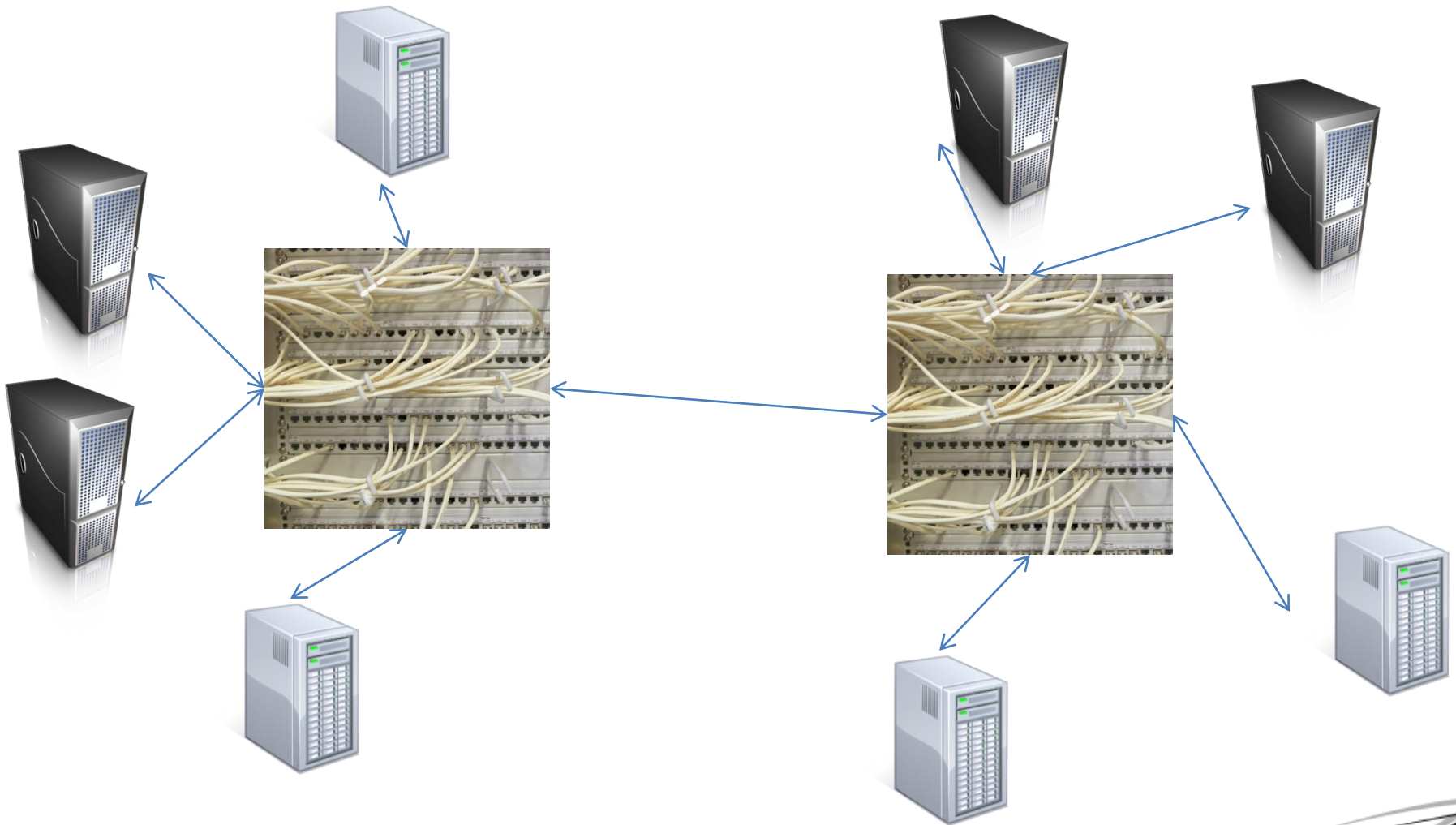
# Parallel on one box

- Threads – shared the same memory
- Processes – special inter-process communication

# Modern cluster architectures

- Special-purpose high-performance clusters (HP, IBM, Sun Cray – hardware & software solutions)
- Commodity hardware + cheap network

# Modern clusters



# Cluster features

- Slow exchange between nodes
- Failures
- Balancing

# Cloud computing

- Cloud application – Internet applications
- Cloud infrastructure – distributed cluster of computers
- Cloud platform – application framework (script language, storage)
- Cloud service – services that provided by cloud computers

# Summary

- Smaller structures can be faster
- In-memory structures for dictionary
- Speller
- Cluster structures



# Q&A